

Around Model Theory Around Free Groups And Around That

Eric Jaligot

CNRS - Lyon

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Formulae

- Language of groups: \cdot , $^{-1}$, 1 .
- Group equations with variables x
- Finite sentences: And, Or, Not, \forall , \exists .

No free variables: sentences

Free variables: $- >$ definable sets

May allow parameters

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Definable sets

Fix a group G $(\cdot, ^{-1}, 1)$.

The truth of a sentence is naturally defined.

$$\varphi(x, a)$$

Definition

The set of tuples g of G such that $\varphi(g, a)$ is true is *definable*.
(by the formula $\varphi(x, a)$ with parameters a)

Quantifier elimination?

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Examples

- **Exemples of sentences.**

Axioms of groups.

Commutativity.

Bounded simplicity.

- **Exemples of definable sets.**

Center. Commutators (but not derived subgroups).

Squares, cubes, etc...

$C(a)$, a^G , translates, etc...

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For Φ a set of sentences (or an *elementary theory* T), the *Φ -elementary class* is the set of groups satisfying all φ in Φ .

- Φ *consistent*: The Φ -elementary class is not empty (*cptness*).
- Groups in the Φ -elementary class: *Models* of Φ .
- Φ *complete*: Consistent + Maximal.
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A theory T has *quantifier elimination* if every formula $\varphi(x)$ is equivalent modulo T to a quantifier-free formula.

Fact (Tarski - Chevalley)

Algebraically closed fields have quantifier elimination.

Fact

Abelian groups eliminate up to boolean combination of cosets.

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CSA-groups

Definition (CSA-groups)

Maximal abelian subgroups are malnormal.

Fact

$\{\text{CSA}\} = \{\text{Centralizers are abelian and selfnormalizing}\}$

In particular: CSA is an elementary class (universal axioms).

$\{\text{Free gps}\} \subseteq \{\text{torsion-free hyperbolic gps}\} \subseteq \{\text{CSA-gps}\}$

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More elementary properties

Universal axioms: For p prime, no elementary abelian p -group of order p^{n+1} .

Fact

The class of CSA-groups with fixed rank of maximal abelian p -subgroup and without involutions is closed under:

- *Free products with 1-malnormal amalgamated subgroup.*
- *HNN-extensions on malnormal separated subgroups.*

Corollary (Ould Houcine)

Existentially closed CSA-group with fixed rank of maximal abelian p -subgroup and without involutions are divisible, with conjugate maximal abelian subgroups, and boundedly simple.

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Definitions

$G_1 \leq G_2$: G_1 is a subgroup of G_2 .

Example

- $F_2 \simeq \langle a, b \mid \mid \rangle \leq \langle a, b, r \mid r^n = a \rangle \simeq F_2$
- $F_2 \simeq \langle a, b \mid \mid \rangle \leq \langle a, b, t \mid a^t = b \rangle \simeq F_2$

The most favorable case:

Definition

$G_1 \preceq G_2$ is an *elementary extension* if $G_1 \leq G_2$ and for every formula $\varphi(x)$ and g_1 in G_1 , $\varphi(g_1)$ true in G_1 implies $\varphi(g_1)$ true in G_2 .

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Chains

Fact (Tarski's test)

$G_1 \preceq G_2$ iff for every formula $\varphi(x, y)$ (y 1-uple) and g_1 in G_1 , if G_2 satisfies $\exists y \varphi(g_1, y)$ then $\varphi(g_1, \gamma)$ for some γ in G_1 .

Fact (Union of chains)

Let $(G_i)_{i < \gamma}$ s.t. $G_i \preceq G_j$ whenever $i \leq j$. Then $G_i \preceq \bigcup_{i < \gamma} G_i$.

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- Existential case.

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Generators and relations

A group G is always given by generators A and relations R .

$$G = \langle A \mid R \rangle$$

Fact (Gromov)

An arbitrarily chosen finitely presented group is hyperbolic with probability almost one.

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Possible relations

Γ a (possibly oriented) irreflexive graph on n vertices

Theorem (Muranov Neman)

For most group words $w(x, y)$, the group

$$\langle a_1, \dots, a_n \mid w(a_i, a_j); \Gamma(a_i, a_j) \rangle$$

is torsion-free hyperbolic and $w(a_i, a_j) = 1$ iff $\Gamma(a_i, a_j)$.

Proof: $C'(1/6)$, hyperbolicity, asphericity. □

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Complexity

Corollary

Fix an arbitrary group word $w(x, y)$. Then for every infinite set of finite graphs Γ_k , there exists a torsion-free hyperbolic group G_k with elements a_1, \dots, a_k s.t.

$$w(a_i, a_j) = 1 \text{ iff } \Gamma(a_i, a_j)$$

Maximal local complexity of {t.f. hyperbolic groups}.

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Approach by model-theoretic complexity of definable sets.
(usually for an elementary class).

Independence graphs: Finite *bipartite* graphs Γ_n on $n + 2^n$ elements coding the powerset of a set of n elements (Vapnik-Chervonenkis).

Definition

A formula $\varphi(x, y)$ has the **Independence Property** relative to a class \mathcal{G} of groups (not nec. elem.) if for each n there exists G_n in \mathcal{G} such that in G_n the definable set defined by $\varphi(x, y)$ induces an independence graph Γ_n .

Ex: $[x, y] = 1$ in the group of permutations of finite support of an infinite set (Zilber - Belegradek - Baldwin-Saxl).

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Ex: Any infinite linear order, with $x \leq y$.

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$$(\text{NOP} \implies \text{NIP})$$

Order Property

Order graphs: Finite *bipartite* graphs Γ_n on $n + n$ elements coding a maximal chain in the powerset of a set of n -elements.

Definition

A formula $\varphi(x, y)$ has the **Order Property** relative to a class \mathcal{G} of groups (not nec. elem.) if for each n there exists G_n in \mathcal{G} such that in G_n the definable set defined by $\varphi(x, y)$ induces an order graph Γ_n .

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Back to {t.f. hyperbolic groups}

- Most group words have the IP relative to the class of torsion-free hyperbolic groups (countably many groups).
- Transfers to existentially closed CSA-groups (2^\perp).
Phenomenon **antipodal** to algebraically closed fields.

What if *finitely* many t.f. hyperbolic groups? *One?*

$\exists x_1 \cdots x_n, y_1, \cdots y_{2^n} \varphi(x, y)$ codes the indep. graph Γ_n

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Elementary classes

Definitions

Φ -elementary class (usually Φ complete theory).

Definition

$\varphi(x, y)$ defines a **stable** set if $\varphi(x, y)$ **does not** have the OP relative to the Φ -elementary class.

It means: there is a uniform bound n for which $\varphi(x, y)$ encodes order graphs Γ_n . \rightarrow **stability index** of φ .

$\forall x_1 \cdots x_{n+1}, y_1 \cdots y_{n+1}$ boolean combination of $\varphi(x_i, y_j)$

Stable sets are closed under boolean combinations and adjunctions of parameters.

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Stable groups

Definition

A group is *stable* if every definable set $\varphi(x, y)$ is stable.

Remark

The stability index of each definable set φ is witnessed by a “ $\forall\varphi$ ” formula in the elementary theory of G .

Sela

Theorem (Sela)

Any torsion-free hyperbolic group is stable.

Quantifier elimination up to $\forall\exists$ sets

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Quotients!

Example (Folklore)

Stable groups may have unstable quotients (for ex. plenty of unstable quotients of free groups).

Example (Meirembekov)

$\langle (x_i), z \mid x_i^3; z^3; [x_i, z]; [x_i, x_j] = z, i < j; [x_i, x_j] = z^2, j > i \rangle$

$Z(G) = \langle z \rangle$ cyclic of order 3

$G/Z(G)$ elementary abelian of order 3 (\aleph_1 -categorical)

$[x_i, x_j] = z$ iff $i < j!$

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Conjectures

Question (Famous in model theory!)

Build new stable groups.

Question

Is the free product of two stable groups still stable?

Corredor's construction

- Start with a free group F_2 .
- Conjugate maximal abelian subgroups by successive *HNN*-extensions, and take the union.
- Repeat countably many times, and take the union.

CSA-group with **conjugate** maximal abelian subgroups.

Similar if one wants to force divisibility.

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Modifications

Lemma

For each n , any t.f. CSA-group G_1 with cyclic centralizers embeds in such a group G_2 in such a way that maximal abelian subgroups of G_1 are conjugate and elements of G_1 has n -th roots.

Proof

HNN-extensions. Then add the n -th root. □

Start with G_1 a free group (stable).

$G_1 \leq G_2 \leq \dots \leq G_n \leq \dots$ where $G_{n-1} \leq G_n$ as in the lemma.

In the union maximal abelian subgroups are conjugate and divisible.

What sets are stable?

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Generic sets

G stable group.

Definition

$X \subseteq_{\text{def}} G$ is *left generic* if finitely many left translates cover G .

Fact

- *Left-genericity is equivalent to right-genericity.*
- *If $X \cup Y$ is generic, then one of X or Y is generic.*

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Connectivity

Definition

G is *connected* if all definable subgroups of finite index are G .

Fact

G is connected iff no partition into two definable generic subsets.

Remark (Poizat)

If X is a definable generic subset of $F = \langle e_n \mid \mid \rangle$, then all but finitely many e_n are in X .

$F = g_1 X \cup \dots \cup g_s X$, e_1, \dots, e_r all generators involved. $e_{r+1} \dots \in X$.

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