

# Catalytic majorization in quantum information theory

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# Majorization

Let  $P_d = \{x \in \mathbf{R}^d, x_i \geq 0, \sum x_i = 1\}$  ;  $P_d \subset P_{d+1} \subset \dots$   
 $x^*$  : decreasing rearrangement of  $x$ .

## Definition

Let  $x, y \in P_d$  ;  $x$  is majorized by  $y$  ( $x \prec y$ ) if

$$\forall k = 1, \dots, d \quad \sum_{i=1}^k x_i^* \leq \sum_{i=1}^k y_i^*$$

$\forall x \in P_d, (\frac{1}{d}, \dots, \frac{1}{d}) \prec x \prec (1, 0, \dots, 0)$ .

## Proposition

$x \prec y \iff \text{There is a bistochastic matrix } B \text{ so that } x = By$ .

For  $y \in P_d$ , let  $S_d(y) = \{x \in P_d \text{ s.t. } x \prec y\}$ .

Then  $S_d(y) = \text{conv}\{(y_{\sigma(1)}, \dots, y_{\sigma(d)}), \sigma \in \mathfrak{S}_d\}$  is a convex polytope.

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# In quantum information theory

Alice and Bob share a quantum system: its states are given by unit vectors in a Hilbert space  $\mathcal{H}_A \otimes \mathcal{H}_B$ .

## Theorem (Nielsen)

Let  $\phi$  and  $\psi$  be two such states and write their Schmidt decomposition

$$\phi = \sum_{i=1}^{d_1} \sqrt{x_i} e_i^A \otimes e_i^B \quad \text{and} \quad \psi = \sum_{i=1}^{d_2} \sqrt{y_i} f_i^A \otimes f_i^B$$

Then Alice and Bob can transform  $\phi$  into  $\psi$  using local quantum operations and classical communication if and only if  $x \prec y$ .

Example : a state is separable if it is of the form  $e^A \otimes e^B$ . Separable states can be transformed only into separable states since no vector majorizes  $(1, 0, \dots, 0)$ .

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# Catalysis and multiple-copy transformations

- If  $x \in P_d, x' \in P_{d'},$  then  $x \otimes x' \in P_{dd'}.$
- Note that if  $x_1 \prec y_1$  and  $x_2 \prec y_2,$  then  $x_1 \otimes x_2 \prec y_1 \otimes y_2$  (tensorize bistochastic matrices).
- The converse is false: it can happen that
  - ①  $x \not\prec y$  but  $\exists z$  so that  $x \otimes z \prec y \otimes z$  (catalysis).
  - ②  $x \not\prec y$  but  $\exists n$  so that  $x^{\otimes n} \prec y^{\otimes n}$  (multiple-copies transformation).
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## Catalysis and multiple-copy transformations (2)

Introduce the following sets which generalize  $S_d(y)$  :

$$C_d(y) = \{x \in P_d \text{ s.t. } \exists k \in \mathbf{N}, \exists z \in P_k, x \otimes z \prec y \otimes z\}$$

$$M_d(y) = \{x \in P_d \text{ s.t. } \exists n \in \mathbf{N}, x^{\otimes n} \prec y^{\otimes n}\}$$

- ①  $C_d(y)$  is convex
- ②  $M_d(y) \subset C_d(y)$  — use  $z = \frac{1}{n}(x^{\otimes(n-1)} \oplus x^{\otimes(n-2)} \otimes y \oplus \cdots \oplus y^{\otimes(n-1)})$
- ③ One may need arbitrarily large  $z$  and  $n$
- ④ These sets are typically not closed

Question (#4 in R.Werner's list of QIT open problems)

How to describe the convex body  $\overline{C_d(y)}$  ? Does  $\overline{C_d(y)} = \overline{M_d(y)}$  ?

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## Necessary conditions

- A function  $f : \mathbf{R}^d \rightarrow \mathbf{R}$  is Schur-convex if  $x \prec y$  implies  $f(x) \leq f(y)$ .
- A function  $f : \mathbf{R}^d \rightarrow \mathbf{R}$  is Schur-concave if  $x \prec y$  implies  $f(x) \geq f(y)$ .
- For  $x \in P_d$  and  $p \in \mathbf{R}$ , let

$$N_p(x) = \sum_{i=1}^d x_i^p.$$

- $N_p$  is Schur-convex for  $p \geq 1$  or  $p \leq 0$  and Schur-concave for  $0 \leq p \leq 1$ . Also,  $N_1(x) = 1$  and  $N_0(x) = d$ .
- $N_p$  is multiplicative :  $N_p(x \otimes z) = N_p(x)N_p(z)$ .
- Therefore  $x \in C_d(y)$  or  $x \in M_d(y)$  imply  $N_p(x) \leq N_p(y)$  or  $N_p(x) \geq N_p(y)$  depending on  $p$ .
- Are there other conditions ?

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# Nielsen's conjecture and our results

## Conjecture (Nielsen)

Let  $y \in P_d$ ,  $y_i > 0$ . Then  $\overline{C_d(y)} = \overline{M_d(y)}$  = the set of  $x \in P_d$  so that

- (A)  $N_p(x) \leq N_p(y)$  for  $p \geq 1$ .
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We can prove partial results

## Theorem (Aubrun+Nechita)

If  $x, y \in P_d$  satisfy (A) and (B), then  $x \in \overline{M_{d+1}(y)}$ ; idem with  $C_{d+1}(y)$ .

## Theorem (Aubrun+Nechita)

If  $x, y \in P_d$  satisfy (A), then  $x \in \overline{\bigcup_{n \geq d} M_n(y)}^{\ell_1}$ ; idem with  $\bigcup C_n(y)$ .

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# From vectors to random variables

- Idea (cf Kuperberg) : to a vector  $x$  associate the random variable

$$V_x \sim \sum_{i=1}^d x_i \delta_{\log x_i}.$$

- Then  $V_{x+y} \stackrel{\perp\!\!\!\perp}{\sim} V_x + V_y$ .
- Rather than majorization on vectors, we will use the more flexible stochastic ordering on random variables

$$X \leq_{st} Y \iff \forall t \in \mathbf{R}, P(X \geq t) \leq P(Y \geq t)$$

Equivalently  $\tilde{X} \leq \tilde{Y}$  on some probability space (Strassen's theorem).

## Lemma

If  $V_x \leq_{st} V_y$ , then  $x \prec y$ .

**Remark :**  $V_x \leq_{st} V_y$  can hold only when  $\text{size}(x) > \text{size}(y)$ . This is why this approach fails to prove the complete conjecture.

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Equivalently  $\tilde{X} \leq \tilde{Y}$  on some probability space (Strassen's theorem).

## Lemma

If  $V_x \leq_{st} V_y$ , then  $x \prec y$ .

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# From vectors to random variables

- Idea (cf Kuperberg) : to a vector  $x$  associate the random variable

$$V_x \sim \sum_{i=1}^d x_i \delta_{\log x_i}.$$

- Then  $V_{x \otimes y} \stackrel{\perp}{\sim} V_x + V_y$ .
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# $N_p(\cdot)$ and Laplace transforms

Recall that  $V_x \sim \sum_{i=1}^d x_i \delta_{\log x_i}$ . Then  $N_{p+1}(x) = \mathbf{E} \exp(pV_x)$ .

Inequalities on  $N_p$  translate into inequalities on the Laplace transforms.

Theorem (Cramér's large deviations theorem)

Let  $X$  be a r.v. and assume  $\Lambda(\lambda) := \log \mathbf{E} e^{\lambda X} < +\infty$ . Set

$$\Lambda^*(t) = \sup_{\lambda \in \mathbb{R}} \lambda t - \Lambda(\lambda).$$

Then for all  $t \in (\min X, \max X)$ , we have

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \mathbf{P}(X_1 + \cdots + X_n \geq tn) = \begin{cases} 0 & \text{if } t \leq \mathbf{E} X \\ -\Lambda^*(t) & \text{if } t \geq \mathbf{E} X \end{cases}$$

where  $(X_i)$  denote i.i.d. copies of  $X$ .

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# Stochastic ordering of i.i.d. sums

## Corollary

Let  $X, Y$  be two random variables and assume that

- ①  $\forall \lambda > 0, \mathbf{E} e^{\lambda X} < \mathbf{E} e^{\lambda Y} < +\infty,$
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- ③  $\mathbf{E} X < \mathbf{E} Y$
- ④  $\max X < \max Y$
- ⑤  $\min X < \min Y$

Then there exists a  $n \geq 1$  so that

$$X_1 + \cdots + X_n \leq_{\text{st}} Y_1 + \cdots + Y_n$$

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## Sketch of proof

- Define

$$f_n(t) := \mathbf{P}(X_1 + \cdots + X_n \geq nt)$$

$$g_n(t) := \mathbf{P}(Y_1 + \cdots + Y_n \geq nt)$$

We want to show that  $f_n \leq g_n$  for  $n$  large enough.

- Cramér's theorem gives  $\lim \frac{1}{n} \log f_n$  on  $[\mathbf{E}X, \max X]$ , and also  $\lim \frac{1}{n} \log(1 - f_n)$  on  $[\min X, \mathbf{E}X]$ ; and similarly for  $g_n$ .
- Because all inequalities were assumed to be strict, we have also a strict inequality between the limits for  $(f_n)$  and  $(g_n)$ . Limit functions are continuous and monotone on a compact set, so the inequality holds uniformly already for some finite  $n$ .

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# Reformulation using majorization

We can reformulate everything in the language of majorization

## Corollary

Let  $x \in P_{d_x}$  and  $y \in P_{d_y}$  with nonzero coordinates. Assume that

- ①  $N_p(x) < N_p(y)$  for  $1 < p < +\infty$ .
- ②  $N_p(x) > N_p(y)$  for  $-\infty < p < 1$ .
- ③  $H(x) < H(y)$  (where  $H(x) = \sum x_i \log x_i$ ).
- ④  $x_{\max} < y_{\max}$ .
- ⑤  $x_{\min} < y_{\min}$ .

Then there exists an integer  $n$  such that  $x^{\otimes n} \prec y^{\otimes n}$  (i.e.  $x \in M_{d_x}(y)$ ).

From this corollary one can prove our two theorems : replace  $x$  by,  
respectively

$$(x_1 - \frac{\varepsilon}{d}, \dots, x_d - \frac{\varepsilon}{d}, \varepsilon) \quad \text{or} \quad (x_1 - \frac{\varepsilon}{d}, \dots, x_d - \frac{\varepsilon}{d}, \frac{\varepsilon}{k}, \dots, \frac{\varepsilon}{k}),$$

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## summary

Recall the main conjecture :

### Conjecture (Nielsen)

Let  $y \in P_d$ ,  $y_i > 0$ . Then  $\overline{C_d(y)} = \overline{M_d(y)}$  = the set of  $x \in P_d$  so that

- (A)  $N_p(x) \leq N_p(y)$  for  $p \geq 1$ .
- (B)  $N_p(x) \geq N_p(y)$  for  $0 \leq p \leq 1$ .
- (C)  $N_p(x) \leq N_p(y)$  for  $p \leq 0$

$$C_d(y) = \{x \in P_d \text{ s.t. } \exists k \in \mathbf{N}, \exists z \in P_k, x \otimes z \prec y \otimes z\}$$

$$M_d(y) = \{x \in P_d \text{ s.t. } \exists n \in \mathbf{N}, x^{\otimes n} \prec y^{\otimes n}\}$$

$$N_p(x) = \sum_{i=1}^d x_i^p.$$

**Update:** this conjecture has been proved by S. Turgut for  $C_d(y)$  few days after I gave this talk. Cf arxiv:0707.0444